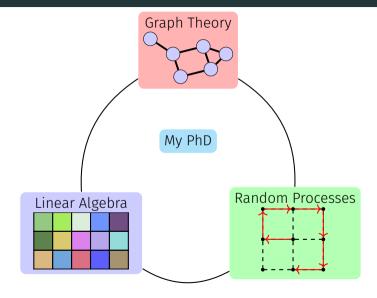
Wilson's Algorithm for Randomized Linear Algebra

Yusuf Yiğit Pilavcı

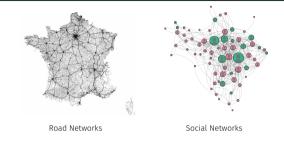
Advisors: Pierre-Olivier Amblard Simon Barthelmé Nicolas Tremblay

WHAT'S INSIDE?

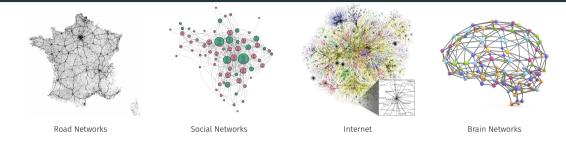


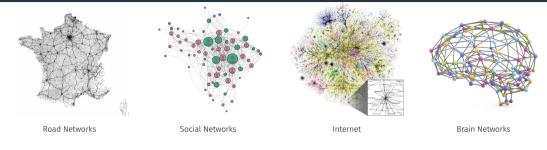


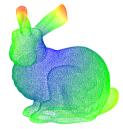
Road Networks



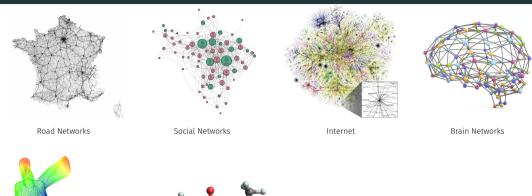








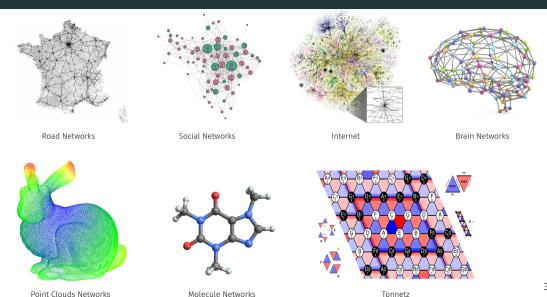
Point Clouds Networks



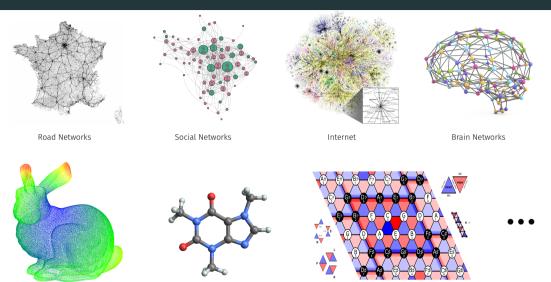




Molecule Networks

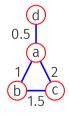


Point Clouds Networks

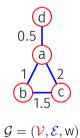


Tonnetz

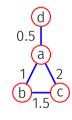
Molecule Networks



 $\mathcal{G} = (\textcolor{red}{\mathcal{V}}, \textcolor{red}{\mathcal{E}}, \textcolor{black}{\mathsf{W}})$



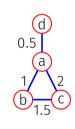
Adjacency matrix W



 $\mathcal{G} = (\mathcal{V}, \mathcal{E}, \mathsf{W})$

Adjacency matrix W

Degree matrix D



 $\mathcal{G} = (\mathcal{V}, \mathcal{E}, \mathsf{W})$

Degree matrix D

Laplacian matrix

$$L = D - W$$

Theory

- · Connectivity Analysis
- Graph Partitioning
- Spanning Trees
- · Random Walks (Loop-Erased)...

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Original Signal:



y:





Figure 3: Median taxi fees paid in drop-off locations in NYC



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Given a graph
$$\mathcal{G} = (\mathcal{V}, \mathcal{E}, w)$$
,
$$\hat{\mathbf{x}} = \arg\min_{\mathbf{x} \in \mathbb{R}^n} q \underbrace{||\mathbf{y} - \mathbf{x}||_2^2}_{\text{Fidelity}} + \underbrace{\mathbf{x}^T L \mathbf{x}}_{\text{Regularization}}, \quad q > 0$$

where L is the graph Laplacian and $\textbf{x}^T L \textbf{x} = \sum\limits_{(i,j) \in \mathcal{E}} w(i,j) (x_i - x_j)^2.$

$$\hat{\mathbf{x}} = K\mathbf{y}$$
 with $K = q(L + qI)^{-1}$

• The explicit solution to this problem is:

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- For large n, iterative methods and polynomial approximations are state-of-the-art.
- For SDD linear systems, there is a growing body of works starting from (Spielman and Teng 2004).

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$$\operatorname{tr}(A) = \sum_{i=1}^{n} \delta_{i}^{\top} A \delta_{i}$$

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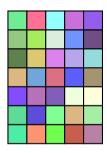
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- \cdot Each uses a quantity called the effective degree of freedom which is equal to $\mathrm{tr}(K)$.

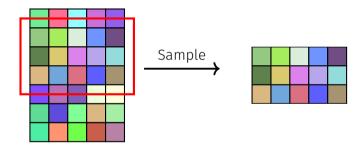
RANDOMIZED LINEAR ALGEBRA

• RLA is a branch of numerical linear algebra developing Monte Carlo methods.



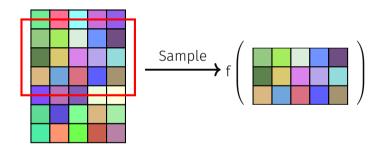
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MAIN THEME

 RLA algorithms for Laplacian-based numerical algebra by using Random Spanning Forests.



OUTLINE

Random Spanning Forests (RSF)

RSF-based Algorithms

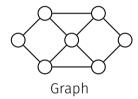
Conclusion

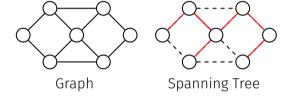
OUTLINE

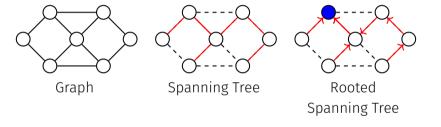
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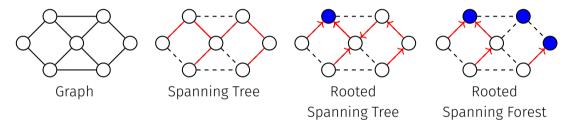
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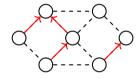






FOREST NOTATIONS

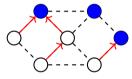
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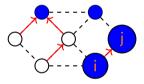
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 \cdot q > 0 changes the expected number of roots.

• The random roots $\rho(\Phi_q)$ is a determinantal point process with a marginal kernel $K=q(L+qI)^{-1}$ (Avena et al. 2018):

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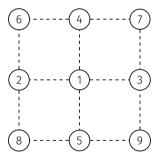
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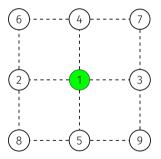
• There is an efficient algorithm to sample RSFs, called Wilson's algorithm (Wilson 1996).

- \cdot Consider an simple random walk on $\mathcal G$ with the transition rule:
 - take a step from i to j with probability $\frac{W(i,j)}{q+d_i}$,
 - · interrupt at any node i with a probability $\frac{q}{q+d_i}$.

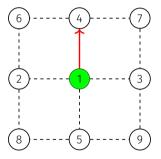
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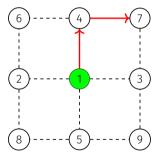
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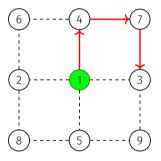
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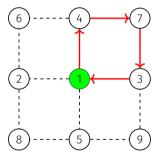
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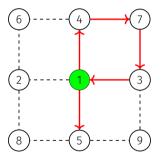
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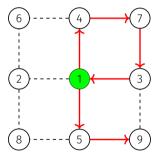
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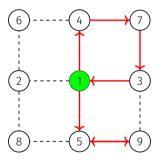
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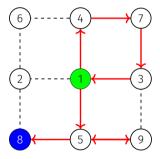
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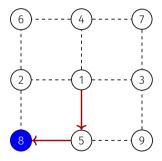
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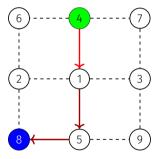
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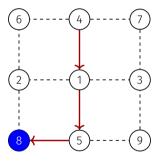
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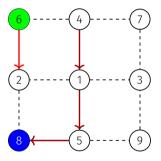
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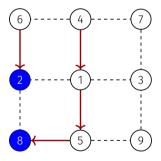
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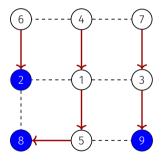
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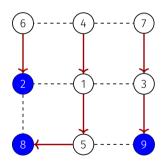
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• The expected number of steps is known:

$$\operatorname{tr}\left[(L+qI)^{-1}(D+qI)\right] \leq \frac{2|\mathcal{E}|}{q} + |\mathcal{V}|.$$

OUTLINE

Random Spanning Forests (RSF)

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Conclusion

MAIN CONTRIBUTIONS

Challenges

- Graph Signal Smoothing
- · Trace Estimation
- Estimating Effective Resistances

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Challenges

 Graph Signal Smoothing Original Signal:



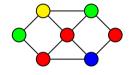


y:

$$\hat{\textbf{x}} = \textbf{K}\textbf{y} :$$



SMOOTHING VIA FORESTS

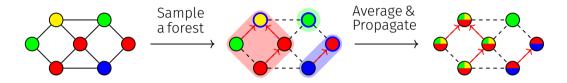


SMOOTHING VIA FORESTS



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SMOOTHING VIA FORESTS



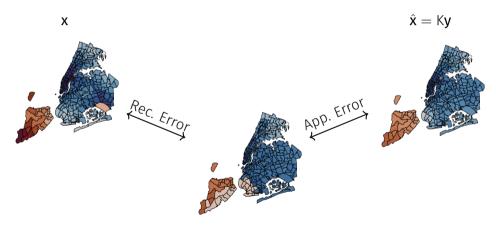
- · Random partitions are sampled via random spanning forests.
- · This yields an unbiased estimator $\bar{\mathbf{x}}$.

COMPARISON WITH STATE OF THE ART

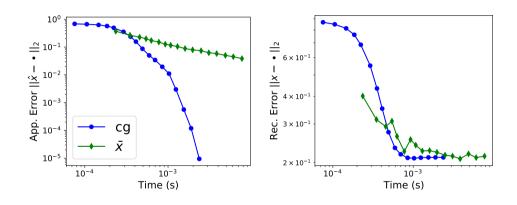
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where $\alpha \in \mathbb{R}$ and $\nabla F(\mathbf{x}_k) = K^{-1}\mathbf{x}_k - \mathbf{y}$.

• We propose to apply the gradient descent update on the previous estimator $\bar{\mathbf{x}}$:

$$\bar{\mathbf{z}} \coloneqq \bar{\mathbf{x}} - \alpha(\mathbf{K}^{-1}\bar{\mathbf{x}} - \mathbf{y})$$

• \bar{z} is unbiased.

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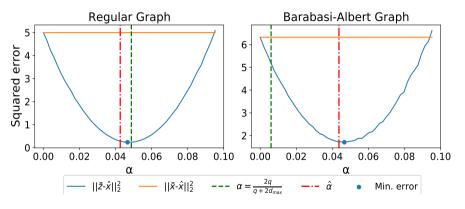
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• One can either choose a value for α from the safe range (e.g. $\alpha = \frac{2q}{q+2d_{max}}$) or estimate from the samples:

$$\hat{\alpha} = \frac{\operatorname{tr}(\widehat{\operatorname{Cov}}(\mathsf{K}^{-1}\bar{\mathbf{x}},\bar{\mathbf{x}}))}{\operatorname{tr}(\widehat{\operatorname{Var}}(\mathsf{K}^{-1}\bar{\mathbf{x}}))}.$$

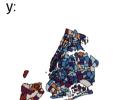
Range of α

• We empirically compare these options of α over a regular and irregular graph:



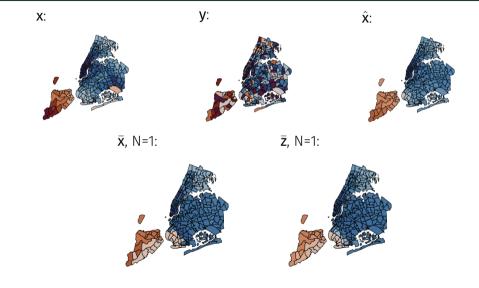
AN ILLUSTRATION







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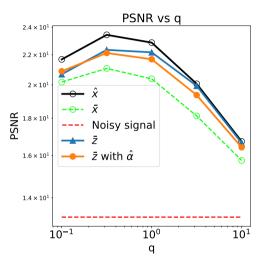


Figure 7: PSNR vs q, N=2

CHALLENGES

- · Graph Signal Smoothing
- · Trace Estimation
- Estimating Effective Resistances

- A famous algorithm for estimating $\operatorname{tr}(K)$ is Hutchinson's estimator:

$$h \coloneqq \frac{1}{N} \sum_{i=1}^{N} \mathbf{a}^{(i)^{\top}} K \mathbf{a}^{(i)}$$

where $\mathbf{a}^{(i)} \in \{-1,1\}^n$ is a random vector with $\mathbb{P}(\mathbf{a}_j^{(i)} = \pm 1) = 1/2$.

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- It is an unbiased estimator of tr(K).
- \cdot The cumbersome computation here is $Ka^{(i)}$ for N vectors.
- It can be done via:
 - · Direct computation via Cholesky decomposition
 - · (Preconditioned) Iterative solvers
 - · Algebraic Multigrid solvers
 - ...

FOREST BASED TRACE ESTIMATOR

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$$\mathsf{s}\coloneqq |
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- One can use this estimator in case of symmetric diagonally dominant matrices instead of the graph Laplacians.

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Variance Reduction via Control Variates

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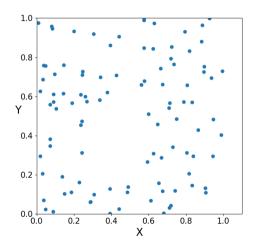
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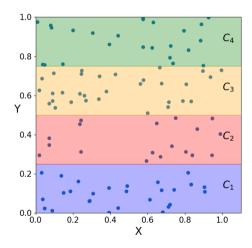
• A safe value of α is $\frac{2q}{q+2d_{max}}$. We also observe that $\frac{2q}{q+2d_{avg}}$ is usually a good estimate of α^* .

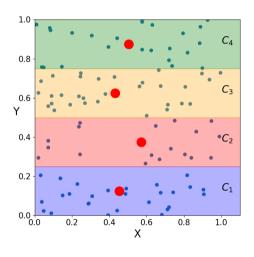
VARIANCE REDUCTION VIA STRATIFICATION

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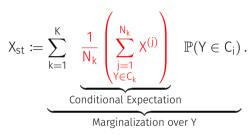


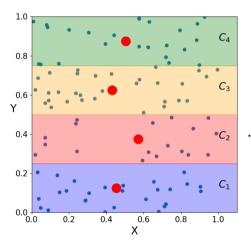
VARIANCE REDUCTION VIA STRATIFICATION





· The stratified estimator is:

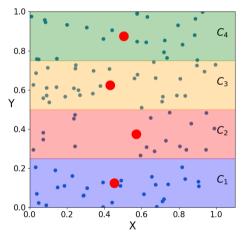




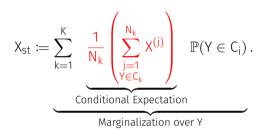
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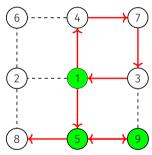
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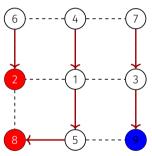
- For certain allocations of N_k 's, one has reduced variance
- We need to have a random variable Y such that:
 - X|Y is easy to sample,
 - $\mathbb{P}(Y \in C_i)$ is accessible.

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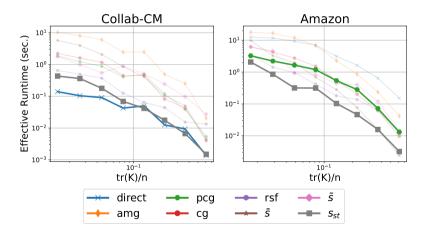


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COMPARISON WITH HUTCHINSON'S ESTIMATOR

 We compare the time needed by the estimators for reaching a certain accuracy.



CHALLENGES

- Graph Signal Smoothing
- Trace Estimation
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· The effective conductance:

$$I_{i,j} \coloneqq \frac{1}{R_{i,i}}$$

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Effective Resistances: What, Why and How?

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- In large scale, they are expensive to compute.

 \cdot The well-known algorithms are Monte Carlo estimators.

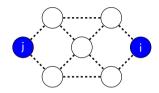
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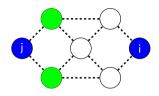
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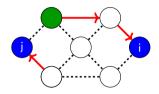
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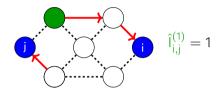
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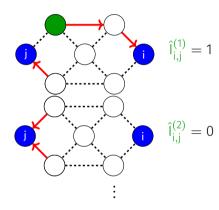
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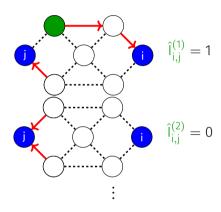






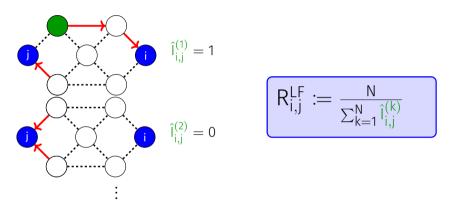






$$\mathsf{R}^{\mathsf{LF}}_{\mathsf{i},\mathsf{j}} \coloneqq \frac{\mathsf{N}}{\sum_{k=1}^{\mathsf{N}} \hat{\mathsf{I}}^{(k)}_{\mathsf{i},\mathsf{j}}}$$

ESTIMATING R_{i,j} VIA LOCAL FORESTS (LF)



· R^{LF} is biased but the bias diminishes faster than the variance.

EXPERIMENTS

• We report the run-time of the local algorithms for approximately the same relative error.

Algorithm Dataset	TP	MC2	LF(ours)
Cora	116	11	2
Citeseer	362	6	1
Pubmed	333	91	12
Collab-CM	82	156	20

Table 1: Runtime (ms) of the local algorithms over benchmark datasets

OUTLINE

Random Spanning Forests (RSF)

RSF-based Algorithms

Conclusion

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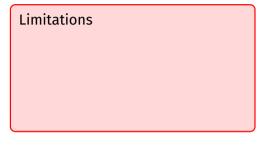
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LIMITATIONS AND OPEN QUESTIONS



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Open Questions

• $\rho_1(\Phi_q)$ is a DPP as well:

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• What happens between $\rho_1(\Phi_q)$ and $\rho(\Phi_q)$?

PUBLICATIONS

Journal

 Yusuf Yiğit Pilavcı, Pierre-Olivier Amblard, Simon Barthelme, and Nicolas Tremblay (2021). "Graph tikhonov regularization and interpolation via random spanning forests". In: IEEE transactions on Signal and Information Processing over Networks 7, pp. 359–374

Conference

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Thanks!



Just a PhD..



RANDOM SPANNING FORESTS

· For fixed subsets $V \subseteq \mathcal{V}$ and $S \subseteq \mathcal{E}$ with |V| = |S|, one has:

$$\det B_{S|V} = \begin{cases} \left(\prod_{(i,j) \in S} w(i,j)\right)^{1/2}, & \text{if S forms a spanning forest rooted in V} \\ 0, & \text{otherwise.} \end{cases}$$

• We can count the spanning forests rooted in $R \subseteq \mathcal{V}$:

$$\forall R \subseteq \mathcal{V}, \quad \det L_{-R} = \sum_{\phi \in \mathcal{F}_R} \prod_{(i,j) \in \mathcal{E}_\phi} w(i,j).$$

• The root probability can be seen as a ratio of counts:

$$\mathbb{P}(r_{\Phi_q}(i)=j) = K_{i,j} = q \frac{(-1)^{i+j}\det(L+qI)_{-i|-j}}{\det(L+qI)} = \frac{|\mathcal{F}^{i\to j}|}{|\mathcal{F}|}$$

LOOP-ERASED RANDOM WALKS

Theorem (Law of LERWs (Marchal 2000)) A loop-erased random walk LE(W) on $\mathcal{G}=(\mathcal{V},\mathcal{E},w)$ that is stopped at the boundary $\Delta\subset\mathcal{V}$ has the following probability distribution:

$$\mathbb{P}(\mathsf{LE}(\mathsf{W}) = \gamma) = \frac{\det \mathsf{L}_{-\Delta \cup \mathsf{S}(\gamma)}}{\det \mathsf{L}_{-\Delta}} \prod_{(i,j) \in \gamma} \mathsf{w}(i,j)$$

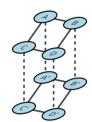
where γ is a fixed path and $s(\gamma)$ denotes the nodes visited in γ .

GRAPH FILTERING VIA DUPLICATED GRAPH

• The product Ky corresponds to a graph filtering with the transfer function:

$$g_{q}(\lambda) = \frac{q}{q + \lambda}$$

· We duplicate the graph and the input $\mathbf{y}_{\mathsf{d}} = \begin{bmatrix} \alpha \mathbf{y} \\ \beta \mathbf{y} \end{bmatrix}$



• The transfer function is paramerized by $\theta = (q_1, q_2, \alpha, \beta)$:

$$f_{\theta}(\lambda) = \frac{\alpha q_1(\lambda + h(0) + q_2) + \beta q_2(h(\lambda))}{(\lambda + h(0) + q_1)(\lambda + h(0) + q_2) - h(\lambda)^2},$$

L₁ Graph Regularization

• Another type of regularization is L₁ regularization:

$$\boldsymbol{x}^{\star} = \operatorname*{argmin}_{\boldsymbol{x} \in \mathbb{R}^{n}} \frac{q}{2} ||\boldsymbol{x} - \boldsymbol{y}||_{2}^{2} + ||\boldsymbol{B}\boldsymbol{x}||_{1}$$

· Alternating direction of multipliers (ADMM) approximates \mathbf{x}^* by:

$$\begin{split} & x_{k+1} = \operatorname*{argmin}_{x \in \mathbb{R}^n} \left(\frac{q}{2} ||x-y||_2^2 + \frac{\rho}{2} ||Bx-z_k+u_k||_2^2 \right) \\ & z_{k+1} = \operatorname*{argmin}_{z \in \mathbb{R}^m} \left(||z||_1 + \frac{\rho}{2} ||Bx_{k+1}-z+u_k||_2^2 \right) \\ & u_{k+1} = u_k + (Bx^{k+1}-z^{k+1}). \end{split}$$

EXTENSION TO SDDS

- · Let $G = U^T \Lambda U = A^{(p)} + A^{(n)} + D^{(1)} + D^{(2)}$ be an symmetric diagonally dominant matrix where $D_{i,i}^{(1)} = \sum_{i \neq j} G_{i,j}$ and $D_{i,i}^{(2)} = G_{i,i} D_{i,i}^{(1)}$.
- · Construct the graph Laplacians:

$$\begin{split} L_1 &= D^{(1)} + A^{(n)} - A^{(p)}/2 = U_1^\top \Lambda_1 U_1 \\ L_2 &= \begin{bmatrix} D^{(1)} + A^{(n)} + D^{(2)}/2 & -D^{(2)}/2 - A^{(p)} \\ -D^{(2)}/2 - A^{(p)} & D^{(1)} + A^{(n)} + D^{(2)}/2 \end{bmatrix} \end{split}$$

- The eigenvectors of L_2 are $U_2 = \begin{bmatrix} U & U_1 \\ -U & U_1 \end{bmatrix}$
- The eigenvalues of L_2 are $\lambda_2 = \lambda_1 \cup \lambda$

CROSS-VALIDATION FOR GTR

• The leave-one-out cross-validation for graph Tikhonov regularization boils down to:

$$LOOCV(q) = \frac{1}{n} \left(\sum_{i=1}^{n} \frac{y_i - \hat{x}_i}{1 - K_{i,i}} \right)^2.$$

• The generalized CV approximation is:

$$GCV(q) = \frac{1}{N} \left(\sum_{i=1}^{n} \frac{y_i - \hat{x}_i}{1 - (\operatorname{tr}(K)/n)} \right)^2.$$

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